



Indra Ganesan

COLLEGE OF ENGINEERING

Approved by AICTE, New Delhi & Affiliated to Anna University, Chennai
Accredited by NAAC with 'B+' Grade, 2(f) & 12B Status Institution by UGC

IG Valley, Madurai Main Road, Manikandam, Tiruchirappalli - 620012

NAAC DOCUMENTS

QUALITY INDICATOR FRAME WORK

CRITERION – 1

CURRICULAR ASPECTS

SUBMITTED BY

IQAC

INTERNAL QUALITY ASSURANCE CELL

INDRA GANESAN COLLEGE OF ENGINEERING





Indra Ganesan

COLLEGE OF ENGINEERING

Madurai Main Road (NH-45B), Manikandam, Tiruchirappalli - 620 012
Approved by AICTE, New Delhi & Affiliated to Anna University, Chennai
NAAC Accredited, 2(F) Status Institution by UGC



Criteria 1	Curricular Aspects	100
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1.1 Curricular Planning and Implementation (20)

1.1.1 The Institution ensures effective curriculum planning and delivery through a well-planned and documented process including Academic calendar and conduct of continuous internal Assessment

Table of Content

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INDRA GANESAN COLLEGE OF ENGINEERING
IG Valley, Manikandam, Tiruchirappalli, Tamil Nadu 620012 , India
(Approved by AICTE, NewDelhi, Affiliated to AnnaUniversity, Chennai-25)

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

PREFACE OF THE COURSE FILE

Batch : 2018-2019

Academic Year : 2018-2019/ODD

Program : M.E. COMPUTER SCIENCE AND ENGINEERING

Year & Semester : 1STYear/1STSemester/'A' Section


Course Code : CP5151

NBA Course Code : C102

Name of the Course : ADVANCED DATA STRUCTURES AND ALGORITHMS

Faculty in-charge : Ms.J.Jenifer/Asst.Prof/CSE


Signature of the Faculty in-charge


Dr. G. Balakrishnan, M.E., Ph.D.,
Principal
Indra Ganesan College of Engineering
IG Valley, Madurai Main Road
Manikandam, Trichy-620 012.


HoD/ CSE

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DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

REVIEW OF COURSE FILE

(To be pasted on the inner side of the file-backside). (#-State Yes/No.)

S.N	Details	Date:	R-I-*	R-II-*&	R-III-*&	R-IV-*&\$	R-V-*&\$@
1.	Preface of the course file		✓				
2.	Vision, Mission, PEOs, POs, PSOs, Blooms taxonomy		✓				
3.	Subject handlers of yesteryears		✓				
4.	Timetable/Workload of the staff - Distribution of teaching load - Roles and Responsibilities		✓				
5.	Syllabus assigned by staff & HoD		✓				
6.	Lecture Schedules assigned by staff & HoD		✓				
7.	Course Committee meeting circular and minutes		✓				
8.	Identification of Curricular gap and Content Beyond the syllabus		✓				
9.	Self-study topics		✓				
10.	Previous AU Question papers		✓				
11.	Unit wise Q & A and Objectivity type questions		✓				
12.	Unit wise course material		✓				
13.	Assignment question paper with sample answers sheets and mark entry			✓	✓	✓	
14.	Tutorial question paper with key and mark entry			✓	✓	✓	
15.	Class test / IA test Q Paper with Key, sample answer papers and mark entry			✓	✓	✓	
16.	IA Test - result analysis - CAP - evidence - root cause analysis.			✓	✓	✓	
17.	Retest - Q paper - Attendance - marks			✓	✓	✓	
18.	AU Web portal entry sheet			✓	✓	✓	
19.	Very poor performance in first two tests - action taken - communication to parents - evidence			✓	✓	✓	
20.	Absence for two tests - action taken - communication to parents - evidence.				✓	✓	
21.	Indiscipline of student reported, if any				✓	✓	
22.	Special class / coaching class / remedial class / attendance - CAP				✓	✓	
23.	Conduct of Seminar, Quizzes - proof						
24.	Content beyond the syllabus - proof			✓	✓	✓	
25.	Student feedback on faculty						
26.	Course end survey						✓
27.	Internal Assessment sheet						✓
28.	AU question paper with students feedback						✓
29.	Discrepancy of the question paper and correspondence, if any						✓
30.	AU result analysis - Detail of far rear students.						✓
31.	AU grade sheet						✓
32.	CO-PO & PSO attainment sheet						✓
	Signature of Course handling faculty						
	Signature of HoD						

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DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

Faculty Time Table

Ms. J. Jenifer AP/CSE								
Day Order	1	2	3	4	5	6	7	8
I			ADA					
II								
III				ADA				
IV							ADA	
V	ADA							
S. Code	Title		Year/Branch		Hours			
CP5151	ADVANCED DATA STRUCTURES AND ALGORITHMS		I/CSEA		4 Hours			
TOTAL- 4hours								


Signature of the Faculty




Hod/CSE

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OBJECTIVES:

- To understand the usage of algorithms in computing.
- To learn and use hierarchical data structures and its operations.
- To learn the usage of graphs and its applications..
- To select and design data structures and algorithms that is appropriate for problems.
- To study about NP Completeness of problems.

UNIT I ROLE OF ALGORITHMS IN COMPUTING 12

Algorithms—Algorithms as a Technology—Insertion Sort—Analyzing Algorithms—Designing Algorithms— Growth of Functions: Asymptotic Notation – Standard Notations and Common Functions- Recurrences: The Substitution Method – The Recursion-Tree Method.

UNIT II HIERARCHICAL DATA STRUCTURES 12

Binary Search Trees: Basics – Querying a Binary search tree – Insertion and Deletion- Red-Black trees: Properties of Red-Black Trees – Rotations – Insertion – Deletion -B-Trees: Definition of Btrees– Basic operations on B-Trees– Deleting a key from a B-Tree- Fibonacci Heaps: structure – Mergeable-heap operations- Decreasing a key and deleting a node-Bounding the maximum degree.

UNIT III GRAPHS 12

Elementary Graph Algorithms: Representations of Graphs – Breadth-First Search – Depth-First Search–Topological Sort – Strongly Connected Components- Minimum Spanning Trees: Growing a Minimum Spanning Tree – Kruskal and Prim- Single-Source Shortest Paths: The Bellman-Ford algorithm – Single-Source Shortest paths in Directed Acyclic Graphs – Dijkstra’s Algorithm; All-Pairs Shortest Paths: Shortest Paths and Matrix Multiplication – The Floyd Warshall Algorithm;

UNIT IV ALGORITHM DESIGN TECHNIQUES 12

Dynamic Programming: Matrix-Chain Multiplication – Elements of Dynamic Programming – Longest Common Subsequence- Greedy Algorithms: An Activity-Selection Problem – Elements of the Greedy Strategy- Huffman Codes.

UNIT V NP COMPLETE AND NP HARD 12

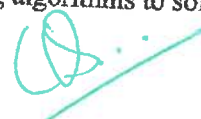
NP-Completeness: Polynomial Time–Polynomial-Time Verification–NP-Completeness and Reducibility – NP-Completeness Proofs – NP-Complete Problems.

Total: 60 PERIODS

OUTCOMES:

At the end of the course, the students should be able to:

- Design data structures and algorithms to solve computing problems.
- Design algorithms using graph structure and various string matching algorithms to solve real-life problems.
- Apply suitable design strategy for problem solving.




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TEXTBOOKS:

1. Alfred Aho, John E. Hopcroft, Jeffrey D. Ullman, — Data Structures and Algorithms I, Pearson Education, Reprint 2006.
2. Robert Sedgewick and Kevin Wayne, — ALGORITHMS I, Fourth Edition, Pearson Education.
3. S. Sridhar, I Design and Analysis of Algorithms I, First Edition, Oxford University Press. 2014.
4. Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest, Clifford Stein, — Introduction to Algorithms I, Third Edition, Prentice-Hall, 2011.

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Signature of the Faculty Incharge


HOD/CSE



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DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

Lecture Schedule

Degree/Program: M.E/CSE

Course Code: CP5151

Course Name: ADVANCED DATA STRUCTURES AND ALGORITHMS

Duration: 2018-2019 Semester: I Section: A Faculty: Ms. J. Jenifer AP/CSE

AIM:

To expose the student's principles of logical data structures and its algorithms

OBJECTIVES:

- To impart knowledge on
- To understand the usage of algorithms in computing.
- To learn and use hierarchical data structures and its operations
- To learn the usage of graphs and its applications.
- To select and design data structures and algorithms that is appropriate for problems.

PREREQUISITES: Graphical Structure theory, Problem Solving algorithms.


AIM:

To expose the student's principles of logical data structures and its algorithms

COURSE OUTCOMES:

After the course, the student should be able to:

CO	Course Outcomes	POs	PSOs
C102.1	Design data structures and algorithms to solve computing problems.	1,2,3,4,5,6,7,8,9,10,11,12	1,2
C102.2	Design algorithms using graph structure.	1,2,3,4,5,6,7,8,9,10,11,12	1,2
C102.3	Design various string matching algorithms to solve real-life problems.	1,2,3,4,5,6,7,8,9,10,11,12	1,2
C102.4	Apply suitable design strategy for problem solving	1,2,3,4,5,6,7,8,9,10,11,12	1,2
C102.5	Develop completeness of algorithm techniques & strategies.	1,2,3,4,5,6,7,8,9,10,11,12	1,2


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S.No	Date	Topics to be Covered	Book
UNIT-I ROLE OF ALGORITHMS IN COMPUTING periods:12			
1	6.7.18	Algorithms	T1/BB
2	7.7.18	AlgorithmsasaTechnology	R2/BB
3	7.7.18	InsertionSort	T1/BB
4	8.7.18	AnalyzingAlgorithms	T3/BB
5	9.7.18	DesigningAlgorithms	R3/BB
6	10.7.18	Growthof Functions	T2/BB
7	15.7.18	AsymptoticNotation	T1/BB
8	16.7.18	StandardNotations	T1/BB
9	17.7.18	CommonFunctions	T1/BB
10	21.7.18	RecurrencesandSubstitutionMethod	
11	21.7.18	TheRecursionandTreeMethod	
UNIT II- HIERARCHICAL DATA STRUCTURES periods:12			
12	28.7.18	BinarySearchTrees.	T1/BB
13	01.8.18	BasicsandQueryingaBinarysearchtree	R2,T1/BB
14	03.8.18	InsertionandDeletion	R2,T1/BB
15	16.8.18	Red-BlacktreesandPropertiesofRed-BlackTrees:	T1/BB
16	20.8.18	Rotations,Insertion,Deletion	R3/BB
17	21.8.18	B-Trees:DefinitionofBtrees	T1/BB
18	24.8.18	BasicoperationsonB-TreesandDeletingakeyfromaB-Tree	T1/BB
19	27.8.18	FibonacciHeaps:structureandMergeable	R1/BB
20	27.8.18	heapoperations	T1/BB
21	28.8.18	Decreasingakeyanddeletinganode	
22	28.8.18	Boundingthemaximumdegree	
UNIT III- GRAPHS Periods:12			
23	30.8.18	ElementaryGraphAlgorithmsandRepresentations ofGraphs	T1/BB
24	31.8.18	BreadthFirst SearchandDepthFirstSearch	T1/BB
25	3.9.18	TopologicalSort	R1/BB
26	5.9.18	StronglyConnectedComponents	T2/BB
27	6.9.18	MinimumSpanningTrees	R1/BB
28	10.9.18	GrowingaMinimumSpanningTree	T3/BB
39	12.9.18	KruskalandPrim- Single	T3/BB
30	24.9.18	SourceShortestPaths: TheBellmanandFordalgorithm	
31	24.9.18	SingleandSourceShortest pathsinDirectedAcyclicGraphsandDijkstra's Algorithm	
32	26.9.18	AllandPairsShortestPaths:ShortestPathsandMatrixMultiplicationand TheFloydWarshall Algorithm	
UNIT IV-ALGORITHM DESIGN TECHNIQUES Target Periods:12			
33	26.9.18	DynamicProgramming	T1/BB
34	28.9.18	Matrix	T1/BB
35	28.9.18	ChainMultiplication	R2/BB
36	2.10.18	ElementsofDynamicProgramming	T1/BB
37	3.10.18	LongestCommonSubsequence	T3/BB
38	4.10.18	GreedyAlgorithms	R3/BB
39	8.10.18	AnActivityandSelectionProblem	R2/BB
40	10.10.18	SelectionProblem	T1/BB
41	10.10.18	ElementsoftheGreedyStrategy	

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42	12.10.18	HuffmanCodes	
UNIT V – NP COMPLETE AND NP HARD		Target Periods:12	
43	12.10.18	NP-Completeness	T1/BB
44	15.10.18	PolynomialTime	T2/BB
45	15.10.18	Polynomial	R1/BB
46	16.10.18	TimeVerification	T3/BB
47	16.10.18	NP-CompletenessandReducability	R3/BB
48	17.10.18	NP-CompletenessProofs	T1/BB
49	17.10.18	NP-Completenessidentification	R2/BB
50	18.10.18	NP-Completenessrootcauses	R1/BB
51	18.10.18	NP-CompletenessProblems	
Content Beyond the Syllabus			
54	27.10.18	REHASHING	Material

TEXTBOOKS:

1. Alfred V. Aho, John E. Hopcroft, Jeffrey D. Ullman, — Data Structures and Algorithms I, Pearson Education, Reprint 2006.
2. Robert Sedgewick and Kevin Wayne, — ALGORITHMS I, Fourth Edition, Pearson Education.
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Signature of the Faculty in-charge

D. G. Balakrishnan
HoD / CSE

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DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

CBS-PROOF

ACADEMIC YEAR: 2018-2019(ODD)

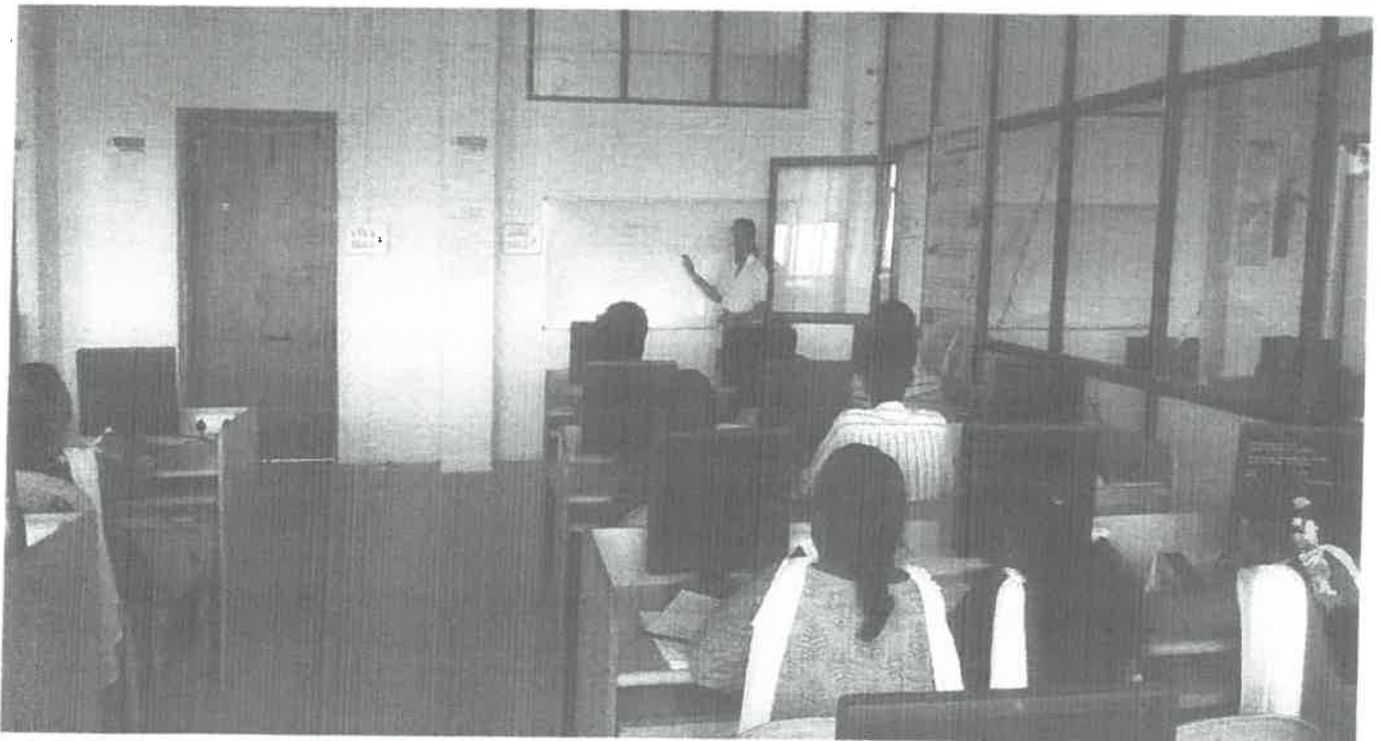
SEM:01


REGULATION:2017

PROGRAM: M.E. CSE

Name of the Faculty: Ms. J. Jenifer

TOPIC: REHASHING SPACE COMPLEXITY




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DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

Assignment Answer Sheet

Name of the Student: *Gwendolyn Rosetta. G*

AU Register Number: *811218405002*

Assignment – 01		Date of Issue:	<i>9.10.2018</i>	Marks	10
Course code	<i>CPS151</i>	Course Title	<i>ADA</i>		
Year	<i>2018</i>	Semester/Section	<i>I/A</i>	Date of Submission:	<i>5.11.2018</i>

Q.No	Questions	CO
1	<i>Analyse Quick Sort, Merge Sort algorithm</i>	<i>2</i>
2	<i>Explain the use of DAC technique for binary search method? Solve its recurrence relation $T(n) = T(n/2) + 1$</i>	<i>3</i>

Mark Allocation

Rubrics	Marks Allocated	Marks obtained
Content Quality	6	<i>4</i>
Presentation Quality	2	<i>1</i>
Timely submission	2	<i>2</i>
Total marks	10	<i>7</i>

Name and Signature of the Faculty Incharge




D. Gudel
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IQAC Academic Audit Form

ACADEMIC YEAR: 2018-2019 ODD SEMESTER

Name of Department : M.E. CSE

Year / Sem / Sec : I / I

No. of Students Registered : 5

Details of Examination : IA Test -1

S.No.	Course Code	List of Reg.No Verified	Course Log Book Verified (Y/N)	Course File Verified (Y/N)	No of students Attended	No of Absentees	No of Failures	Pass %	Remarks
1.	MA5160	811218405001, 811218405004	Y	Y	4	1	-	98%	-
2.	CP5151	811218405001, 811218405004	Y	Y	4	1	-	98%	-
3.	CP5152	811218405001, 811218405004	Y	Y	4	1	-	98%	-
4.	CP5153	811218405001, 811218405004	Y	Y	4	1	-	98%	-
5.	CP5154	811218405001, 811218405004	Y	Y	4	-	-	100%	-
6.	CP5191	811218405001, 811218405004	Y	Y	4	-	-	100%	-

Verified by

External Member Name and Signature:

Internal Member Name and Signature:

Overall Remarks:

D. L. Fudd
HoD/ CSE

IQAC Co-ordinator

Principal

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	<pre>current = circularQueue.front.next // Start from the next node of the front while current is not equal to circularQueue.front: count = count + 1 current = current.next return count.</pre> <p>This algorithm assumes that the circular queue has a structure where each node has a 'next' pointer that points to the next node in the circular sequence, and the last node's 'next' pointer points back to the front of the queue. The 'front' points to the first node in the queue. The algorithm traverses the circular queue, counting nodes until it reaches the front again. If the queue is empty, it returns 0.</p>		
6	<p>Define the terms: Infix, postfix and prefix. In the context of expressions:</p> <ol style="list-style-type: none"> Infix: Infix notation is the standard way of writing mathematical expressions where operators are placed between operands. For example, in the infix expression "a + b", the operator '+' is placed between the operands 'a' and 'b'. Postfix (or Reverse Polish Notation - RPN): Postfix notation is a way of writing mathematical expressions in which the operators come after their operands. For example, in postfix notation, "a + b" is written as "a b +". Evaluation of postfix expressions is often done using a stack data structure. Prefix (or Polish Notation): Prefix notation is a way of writing mathematical expressions in which the operators precede their operands. For example, in prefix notation, "a + b" is written as "+ a b". Like postfix notation, the evaluation of prefix expressions is commonly done using a stack. 	2	2
7	<p>Define Linked list and give its applications. A linked list is a linear data structure where elements are stored in nodes, and each node points to the next node in the sequence, forming a chain-like structure.</p> <p>Applications of Linked Lists:</p> <ol style="list-style-type: none"> Dynamic Memory Allocation: Linked lists allow dynamic memory allocation, making them useful for managing memory efficiently. Implementation of Data Structures: Linked lists are used to implement various data structures like stacks, queues, and symbol tables. Undo Functionality: Linked lists can be employed to implement undo functionality in applications where a history of operations needs to be maintained. File Systems: Linked lists are used in file systems to keep track of the blocks of a file scattered across the storage. Representing Sparse Data: Linked lists are suitable for representing sparse data structures where most elements are empty, saving memory. Polynomial Manipulation: Linked lists are used to represent and manipulate polynomials efficiently. Graph Algorithms: Linked lists can be utilized to represent adjacency lists in graph-related algorithms. 	2	2
8	<p>Define Full binary tree(or) Complete binary tree. A full binary tree is a binary tree in which every node has either 0 or 2 children. In other words, every level of the tree is completely filled, except possibly for the last level, which is filled from left to right. A complete binary tree is a special case of a binary tree where all levels are filled except possibly the last, and all nodes are as left as possible.</p>	2	2
9	<p>Define expression tree. An expression tree is a binary tree representation of an arithmetic expression, where the leaves are operands and the internal nodes are operators. The tree's structure reflects the hierarchy and order of operations in the original expression, making it a useful data structure for evaluating and manipulating mathematical expressions.</p>	2	2
10	<p>Traverse the given tree using Inorder, Preorder and Postorder traversals. Here are brief descriptions of the three tree traversal methods:</p> <ol style="list-style-type: none"> Inorder Traversal: <ul style="list-style-type: none"> Traverse the left subtree in inorder. Visit the root node. Traverse the right subtree in inorder. 	3	2

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```

new_node = Node(value)
if head is None:
    head = tail = new_node
else:
    new_node.prev = tail
    tail.next = new_node
    tail = new_node

```

```

# Deletion from the End
def delete_at_end():
    global head, tail
    if head is None:
        print("Underflow: List is empty.")
    elif head == tail:
        head = tail = None
    else:
        tail = tail.prev
        tail.next = None

```

```

# Display Doubly Linked List
def display():
    current = head
    while current:
        print(current.data, end=" <-> ")
        current = current.next
    print("NULL")

```

```

# Example Usage
create_doubly_linked_list()

```

```

insert_at_end(1)
insert_at_end(2)
insert_at_end(3)

```

```

print("Doubly Linked List after insertion:")
display()

```

```

delete_at_end()

```

```

print("Doubly Linked List after deletion:")
display().

```

This example demonstrates the creation of a doubly linked list, insertion at the end, and deletion from the end. The display function shows the state of the doubly linked list after each operation.

OR

12b **Construct a dequeue data structure in which the following operations to be implemented.**

A deque (double-ended queue) is a data structure that allows elements to be added or removed from both ends. Here's an example implementation in Python using a doubly linked list:

```

class Node:
    def __init__(self, data):
        self.data = data
        self.next = None
        self.prev = None

```

```

class Deque:
    def __init__(self):

```



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2

2

```

new_node.next = head
head = new_node

# Insert elements at the end
temp = head
while temp.next is not None:
    temp = temp.next

new_node = Node(5)
temp.next = new_node

# Delete from the beginning
head = head.next

# Traversal
temp = head
while temp is not None:
    print(temp.data, end=" ")
    temp = temp.next.

```

This example demonstrates the basic operations of inserting at the beginning and end, deleting from the beginning, and traversing a singly linked list.

12a

Explain creation, insertion and deletion of doubly linked list with example.

2

2

Let's go through the creation, insertion, and deletion of nodes in a doubly linked list along with an example.

1. Creation of Doubly Linked List:

- Initialize pointers head and tail to NULL.
- For each element to be inserted, create a new node.
- If the list is empty, set both head and tail to the new node.
- Otherwise, update pointers to maintain the doubly linked structure.

2. Insertion in Doubly Linked List:

- Insertion can occur at the beginning, end, or at a specific position.
- For simplicity, let's consider inserting at the end.
- Create a new node.
- If the list is empty, set both head and tail to the new node.
- Otherwise, update pointers accordingly.

3. Deletion in Doubly Linked List:

- Deletion can occur at the beginning, end, or at a specific position.
- For simplicity, let's consider deleting at the end.
- If the list is empty, indicate underflow.
- Otherwise, update pointers accordingly.

Here's an example in Python:

```

class Node:
    def __init__(self, data):
        self.data = data
        self.next = None
        self.prev = None

# Creation of Doubly Linked List
def create_doubly_linked_list():
    global head, tail
    head = tail = None

# Insertion at the End
def insert_at_end(value):
    global head, tail

```

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11b

Explain the algorithm for implementing Singly Linked list.

2

2

Certainly! Implementing a singly linked list involves creating nodes and linking them in a forward direction. Here's a basic algorithm for the operations on a singly linked list:

Algorithm for Singly Linked List Operations:

1. ****Define Node structure:****
 - Each node contains a data element and a reference (or pointer) to the next node.
2. ****Initialization:****
 - Set the head pointer to null initially, indicating an empty list.
3. ****Insertion at the beginning (InsertFront):****
 - Create a new node with the given data.
 - Set the next pointer of the new node to the current head.
 - Update the head pointer to the new node.
4. ****Insertion at the end (InsertEnd):****
 - Create a new node with the given data.
 - Traverse the list until the last node.
 - Set the next pointer of the last node to the new node.
5. ****Deletion from the beginning (DeleteFront):****
 - If the list is empty, indicate underflow.
 - Otherwise, update the head pointer to the next node.
6. ****Traversal:****
 - Start from the head.
 - While the current node is not null, process the data and move to the next node.
7. ****Search (optional):****
 - Start from the head.
 - While the current node is not null, check if the data matches the target.
 - Return true if found, false otherwise.
8. ****Deletion by value (optional):****
 - Start from the head.
 - Keep track of the current and previous nodes.
 - If the target value is found, update the next pointer of the previous node to skip the current node.

Note: Always handle edge cases such as an empty list, and manage pointers carefully to avoid memory leaks.

Example:

```
python
class Node:
    def __init__(self, data):
        self.data = data
        self.next = None
```

```
# Initialize an empty linked list
head = None
```

```
# Insert elements at the beginning
new_node = Node(3)
new_node.next = head
head = new_node
```

```
new_node = Node(2)
```

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	<p>2. Preorder Traversal:</p> <ul style="list-style-type: none"> • Visit the root node. • Traverse the left subtree in preorder. • Traverse the right subtree in preorder. <p>3. Postorder Traversal:</p> <ul style="list-style-type: none"> • Traverse the left subtree in postorder. • Traverse the right subtree in postorder. • Visit the root node. <p>These traversal methods can be applied to any tree, and the order in which nodes are visited defines the respective traversal. The exact sequence of nodes visited will depend on the specific structure of the tree.</p>		
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PART B
(Answer all the Questions 2x10=20 Marks)

11a	<p>Write the algorithm for performing operations in a stack. Trace your algorithm with suitable Example.</p> <p>Certainly! Let's consider a simple algorithm for performing operations (push and pop) in a stack and trace it with an example:</p> <p>Algorithm for Stack Operations:</p> <ol style="list-style-type: none"> 1. Initialize an empty stack. 2. Push (Insert): <ul style="list-style-type: none"> - Check if the stack is full (overflow). - If not full, insert the element onto the top of the stack. 3. Pop (Delete): <ul style="list-style-type: none"> - Check if the stack is empty (underflow). - If not empty, remove the element from the top of the stack. <p>Example:</p> <p>Let's perform operations on a stack with a maximum capacity of 5:</p> <ol style="list-style-type: none"> 1. Initialize an empty stack: `Stack = []` 2. Push elements onto the stack: <ul style="list-style-type: none"> - Push 1: `Stack = [1]` - Push 2: `Stack = [1, 2]` - Push 3: `Stack = [1, 2, 3]` 3. Pop elements from the stack: <ul style="list-style-type: none"> - Pop: `Stack = [1, 2]` - Pop: `Stack = [1]` 4. Push more elements: <ul style="list-style-type: none"> - Push 4: `Stack = [1, 4]` - Push 5: `Stack = [1, 4, 5]` 5. Pop again: <ul style="list-style-type: none"> - Pop: `Stack = [1, 4]` <p>This algorithm ensures that the stack is not accessed or modified in case of underflow or overflow, respectively. It follows the Last-In-First-Out (LIFO) principle, where the last element pushed is the first one to be popped.</p>	2	2
-----	---	---	---

OR



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subtree. The algorithm modifies the tree and returns the root of the modified BST.

OR

13b **Construct a minimum spanning tree using Kruskal's algorithm with your own example.**

2

3

Kruskal's Algorithm for Minimum Spanning Tree (MST):

Kruskal's algorithm is a greedy algorithm that finds a minimum spanning tree for a connected, undirected graph. It works by sorting the edges in ascending order of their weights and adding edges to the MST while avoiding the formation of cycles.

Let's consider a simple graph and apply Kruskal's algorithm step by step:

Graph:

A - 2 - B

| |
1 3

| |
C - 4 - D

Steps:

Sort Edges by Weight:

- (A, B): 2
- (C, D): 4
- (A, C): 1
- (B, D): 3

Initialize Empty MST:

- MST = {}

Add Edges to MST:

- Add (A, C): 1
 - MST = {(A, C)}
- Add (A, B): 2
 - MST = {(A, C), (A, B)}
- Add (B, D): 3
 - MST = {(A, C), (A, B), (B, D)}

The final Minimum Spanning Tree (MST) contains the edges {(A, C), (A, B), (B, D)} with a total weight of 6.

Note: Kruskal's algorithm can handle graphs with weighted edges, and in case of ties in edge weights, the choice of which edge to add first is arbitrary. The key is to avoid creating cycles in the process.



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Course Faculty



HoD

```

self.front = None
self.rear = None

def is_empty(self):
    return self.front is None

def add_front(self, data):
    new_node = Node(data)
    if self.is_empty():
        self.front = self.rear = new_node
    else:
        new_node.next = self.front
        self.front.prev = new_node
        self.front = new_node

def add_rear(self, data):
    new_node = Node(data)
    if self.is_empty():
        self.front = self.rear = new_node
    else:
        new_node.prev = self.rear
        self.rear.next = new_node
        self.rear = new_node

def remove_front(self):
    if self.is_empty():
        return None
    data = self.front.data
    if self.front == self.rear:
        self.front = self.rear = None
    else:
        self.front = self.front.next
        self.front.prev = None
    return data

def remove_rear(self):
    if self.is_empty():
        return None
    data = self.rear.data
    if self.front == self.rear:
        self.front = self.rear = None
    else:
        self.rear = self.rear.prev
        self.rear.next = None
    return data

def display(self):
    current = self.front
    while current:
        print(current.data, end=" ")
        current = current.next
    print()

# Example Usage:
deque = Deque()

deque.add_front(1)
deque.add_front(2)
deque.add_rear(3)
deque.add_rear(4)

deque.display() # Output: 2 1 3 4

```

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```
print("Removed front:", deque.remove_front()) # Output: 2
print("Removed rear:", deque.remove_rear()) # Output: 4

deque.display() # Output: 1 3
```

In this example, **add front** and **remove front** operations manipulate the front of the deque, while **add rear** and **remove rear** operations manipulate the rear. The doubly linked list structure allows easy addition and removal of elements from both ends of the deque.

PART C
(Answer all the Questions 1x10=10 Marks)

13a

What is Binary search tree? Write an algorithm to add a node into a binary search tree.

2

2

A Binary Search Tree (BST) is a binary tree data structure in which each node has at most two children, referred to as the left child and the right child. For every node, all elements in its left subtree are less than the node, and all elements in its right subtree are greater than the node.

Algorithm to Add a Node to a Binary Search Tree:

Algorithm: InsertNode

Input: Root of the BST, Value to be inserted

1. If the root is null, create a new node with the given value and set it as the root.
2. Otherwise:
 - a. If the value to be inserted is less than the root's value:
 - Recursively call InsertNode on the left subtree.
 - b. If the value to be inserted is greater than the root's value:
 - Recursively call InsertNode on the right subtree.
3. Return the modified root.

Example in Python:

```
class TreeNode:
    def __init__(self, key):
        self.left = None
        self.right = None
        self.val = key

def insert_node(root, key):
    if root is None:
        return TreeNode(key)
    else:
        if key < root.val:
            root.left = insert_node(root.left, key)
        elif key > root.val:
            root.right = insert_node(root.right, key)
    return root
```

Example usage:

```
root = None
keys = [50, 30, 20, 40, 70, 60, 80]
```

```
for key in keys:
    root = insert_node(root, key).
```

In this example, the **insert node** algorithm recursively traverses the BST to find the appropriate position for the new node based on the comparison of values. If the value is less than the current node, it goes to the left subtree; if it is greater, it goes to the right

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Internal Assessment Test Answer Book

Name	Gwendolyn Rosetta . G		Year/Semester/Section	I/I	
Batch No.	811218405002	Date/Session	8.9.2018	Department	M.E.CSE
Course code	CP5151	Course Title	Advance Data Structures & Algorithms		
Internal Assessment Test	IAT1	<input checked="" type="checkbox"/> IAT2	<input type="checkbox"/> IAT3	Model	
Name and Signature of the Invigilator with date					

Instruction to the Student: Put tick mark to the question attended in the column against question.							
PartA			PartB/PartC				TotalMarks
Q.No.	✓	Marks	Q.NO.	✓	a	b	
					Marks	Marks	
1		2	11		10		10
2		2	12		6		6
3		1	13			6	6
4		0	14				
5		2	15				
6		2	16				
7		2				Total	22
8		2	40			 Name and Signature of the Examiner with date	
9		2					
10		2					
Total		18	Grand Total				

To be filled by the examiner							
Course Outcomes	1	2	3	4	5	6	Total
Marks allotted							
Marks Obtained							
IQAC Audit-Remarks  Dr. G. Balakrishnan, M.E., Ph.D., Principal Indra Ganesan College of Engineering IG Valley, Madurai Main Road Manikandam, Trichy-620 012.							Name and Signature of the IQAC member



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DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING
ACADEMIC YEAR 2022 - 2023 (ODD SEMESTER)

STUDENTS MARK STATEMENT- CO BASED
INTERNAL ASSESSMENT TEST-1

SUBJECT CODE & TITLE: CP5151 & ADVANCED DATA STRUCTURES AND
ALGORITHMS

YEAR/SEM: I/I

MONTH & YEAR:

S.NO	REG NO	STUDENT NAME	COX (32)	COX (18)	TOTAL (50)	TOTAL (100)
1.	811218405001	Abwini . M	30	15	45	90
2.	811218405002	Grasendayan Rosetta . G	25	10	35	70
3.	811218405003	Harish . V	20	9	29	58
4.	811218405004	Nirmala . N	24	12	36	72


MARKS RANGE:


<20	20-30	31-40	41-50	51-60	61-70	71-80	81-90	91-100
				1	1	1	1	

Total No. of Candidates Present	4
Total No. of Candidates Absent	0
Total No. of Students Pass	4
Total No. of Students Fail	0
Percentage of Pass	100 %

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